

Secret Sharing

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The cryptographic trust problem

- ◆ In secure systems, an authority is often required to
 - protect some secret, and
 - employ the secret in a manner defined by the system

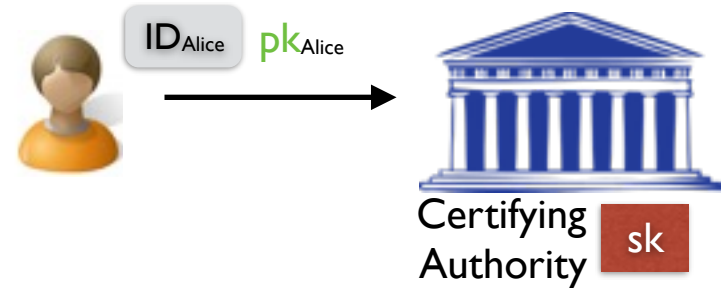
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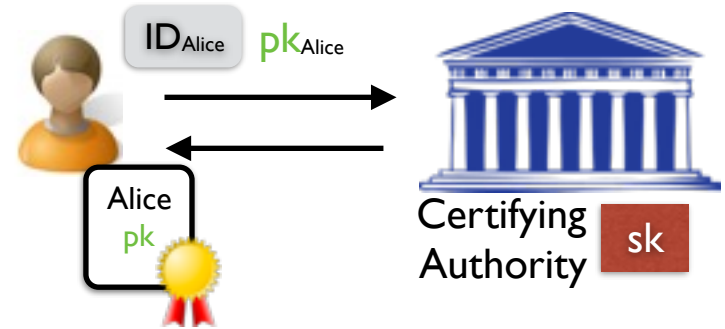
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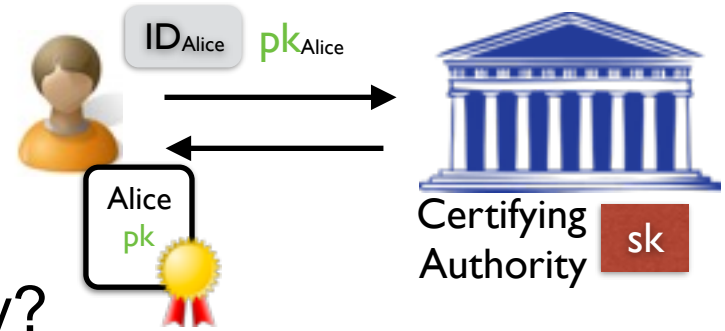
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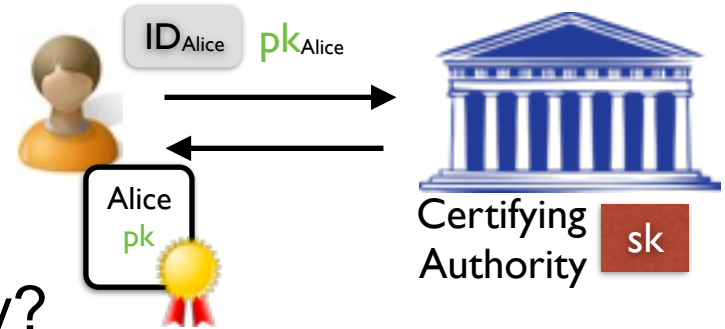
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- ◆ In secure systems, an authority is often required to
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- ◆ How to realize such a trusted party?



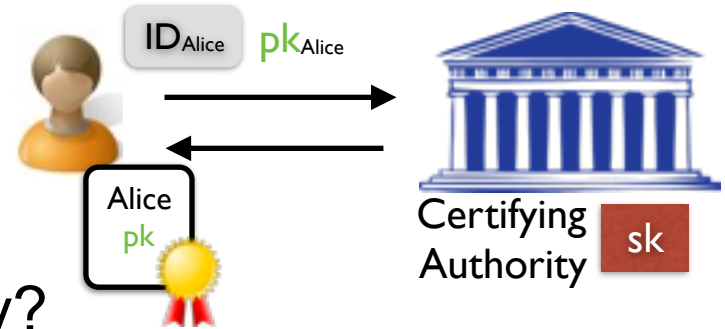
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 - Find such an individual/group/organization, and justify their trustworthiness



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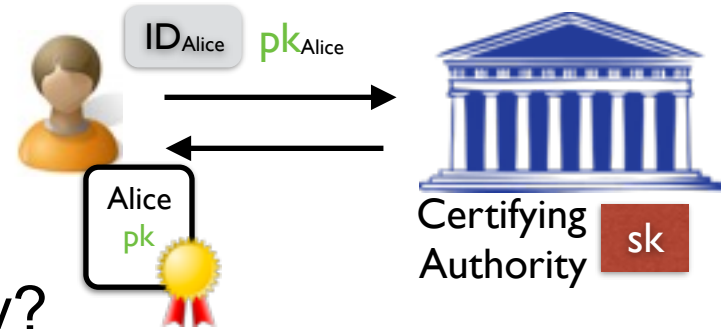


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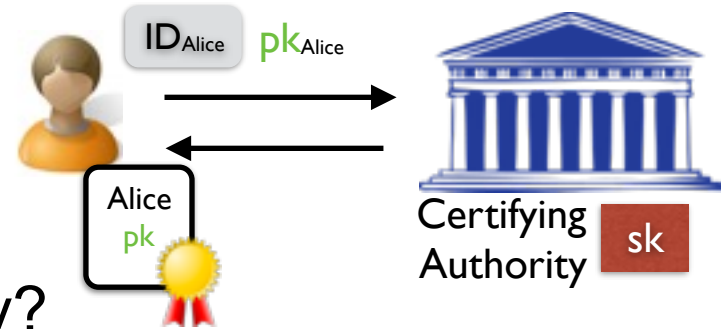
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- Use multi-party computation (MPC) or a distributed authority



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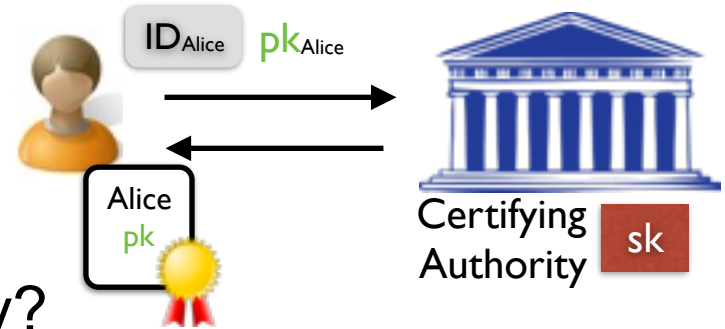


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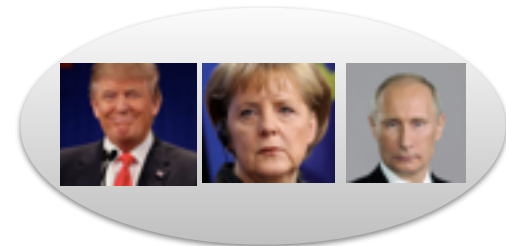


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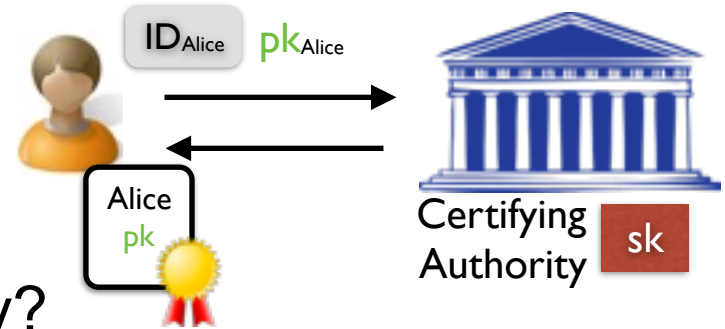


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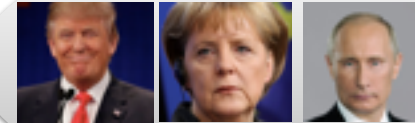


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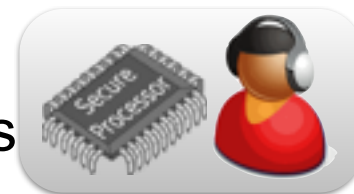
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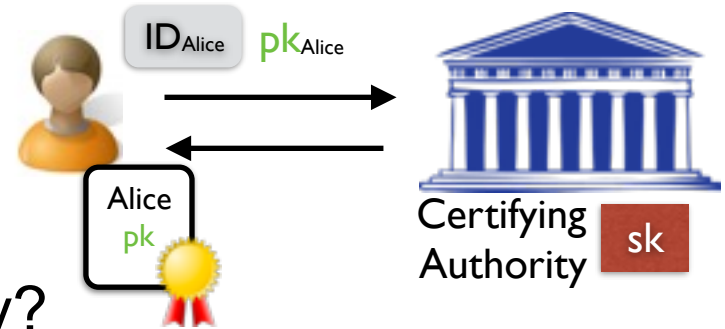


- Employ a trusted hardware module that performs cryptographic operations



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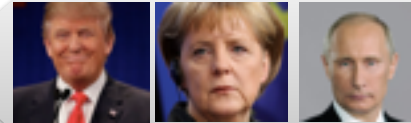


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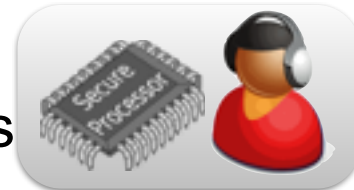
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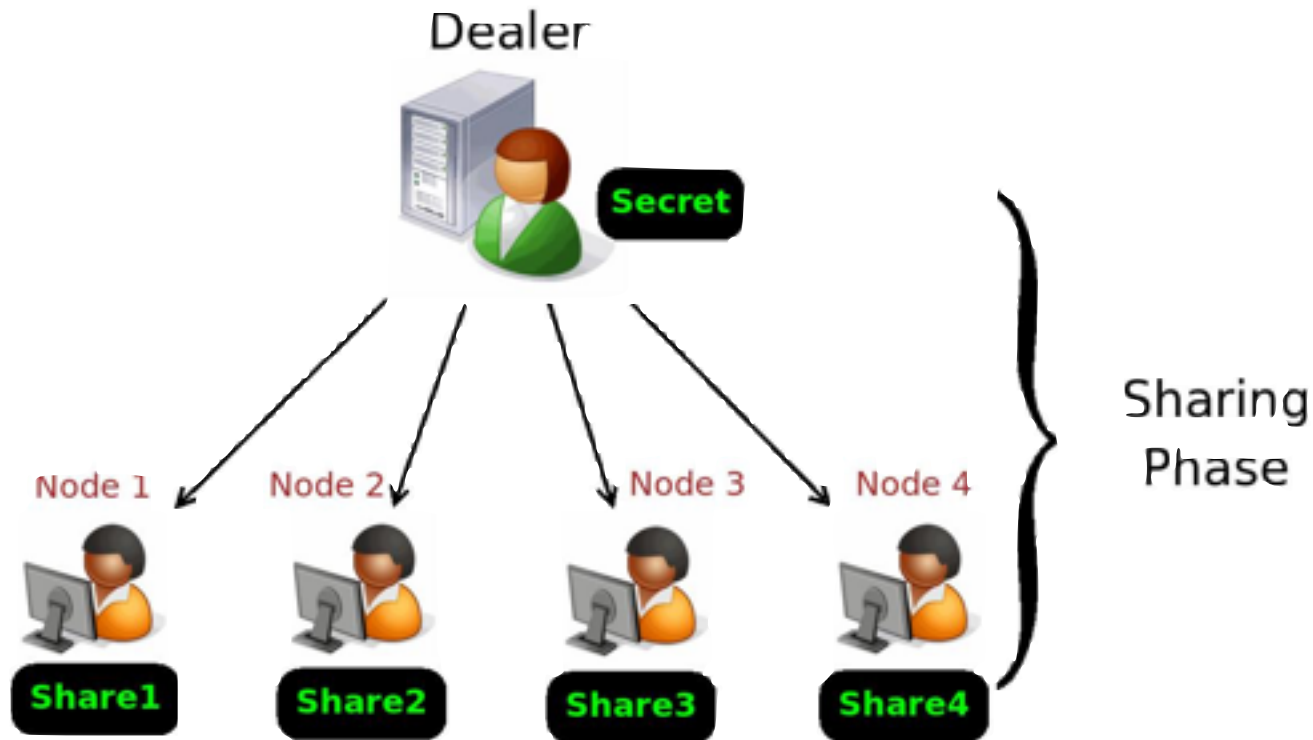


Secret Sharing

- ◆ The notion of secret sharing was introduced independently by Shamir and Blakley in 1979
- ◆ For integers n and t such that $n > t \geq 0$, an (n,t) -secret sharing (SS) scheme is a method used by a dealer D
 - to share a secret s among a set of n parties (the sharing phase) in such a way that
 - in the reconstruction phase any subset of $t+1$ or more parties can compute s , but subsets of size t or fewer cannot
- ◆ What is the relation between n and t
 - against a passive adversary controlling any t parties?
 - against an active adversary controlling any t parties?

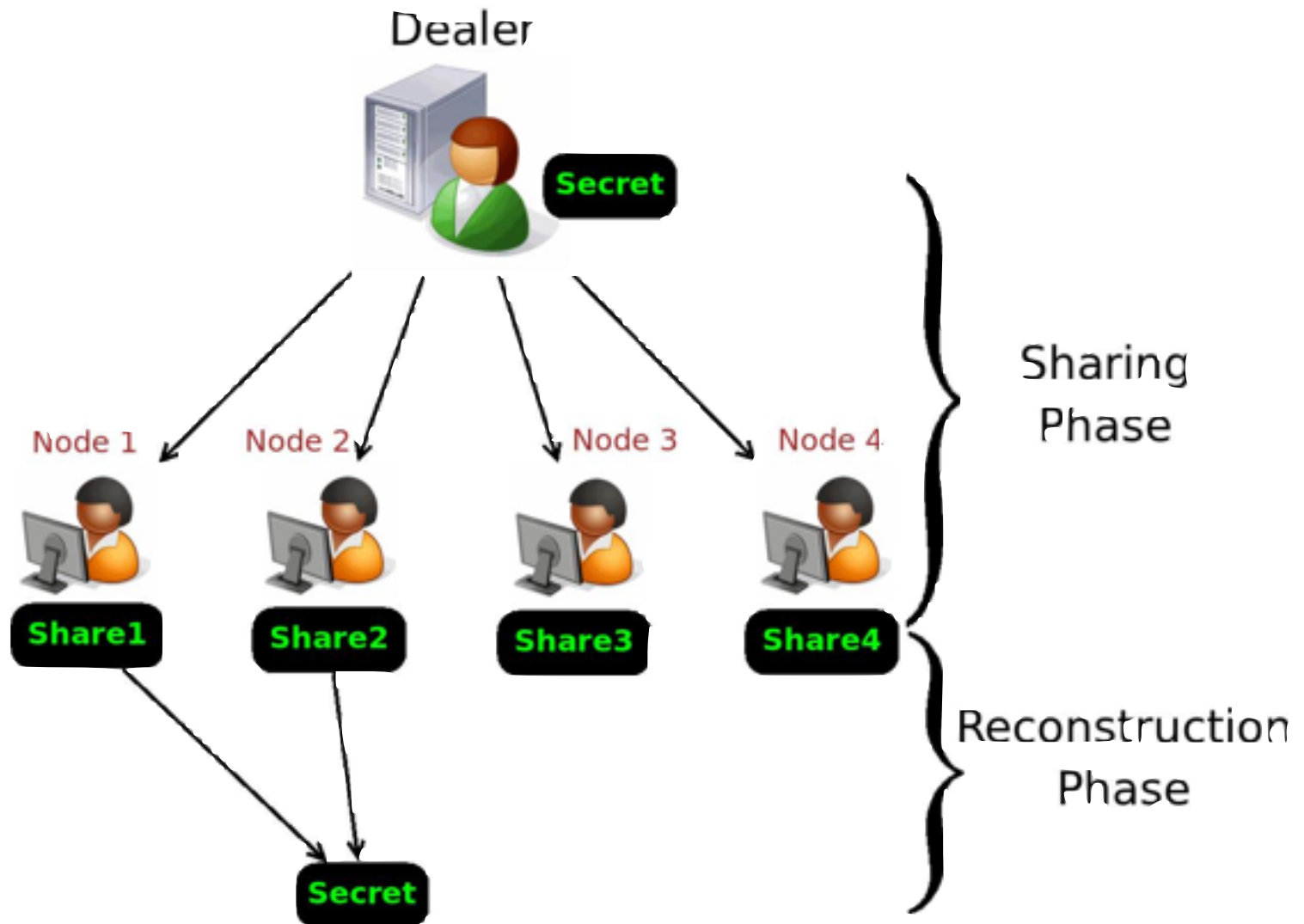
Secret Sharing Basics

◆ Sharing Phase



Secret Sharing Basics

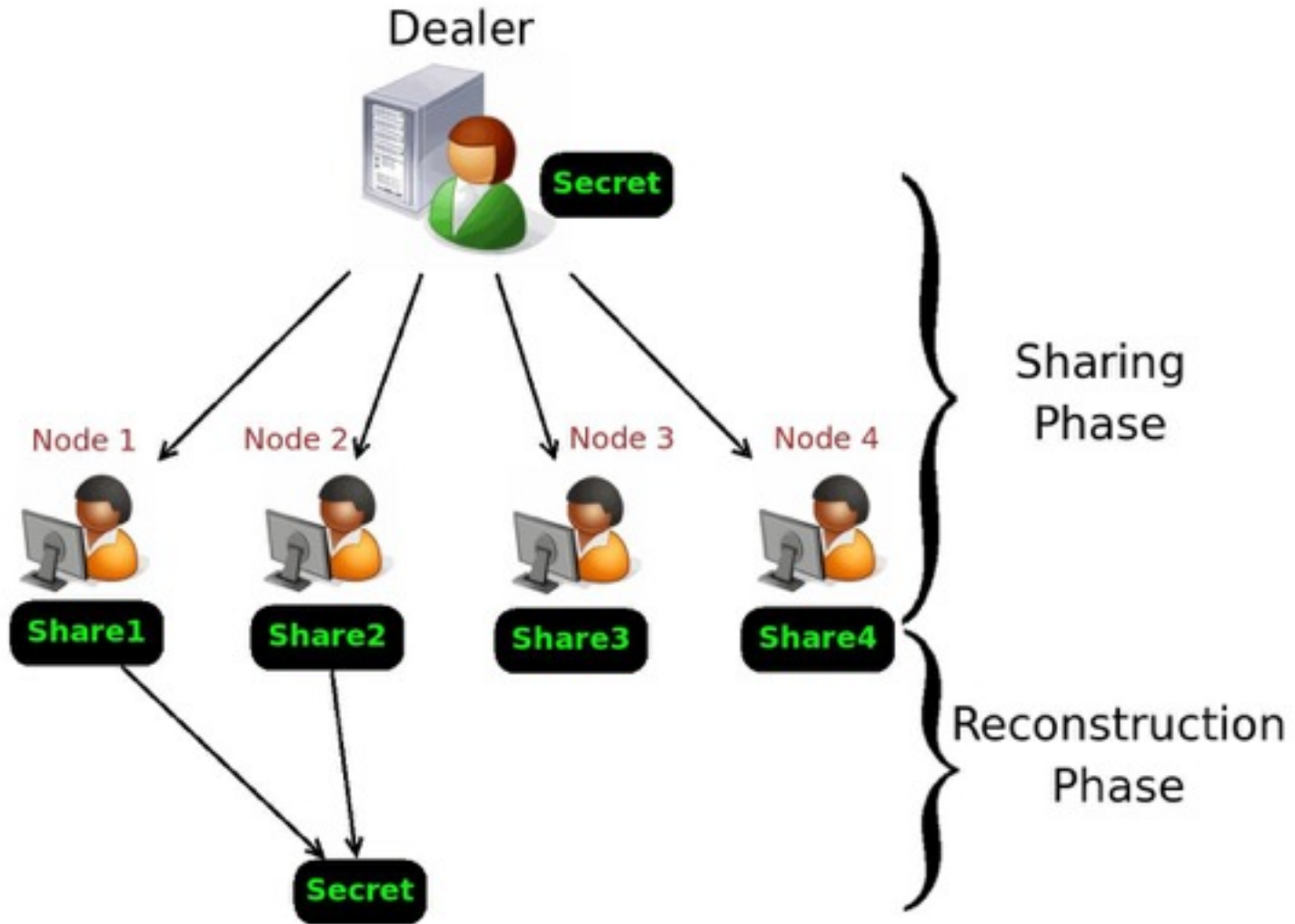
◆ Reconstruction Phase



Shamir Secret Sharing

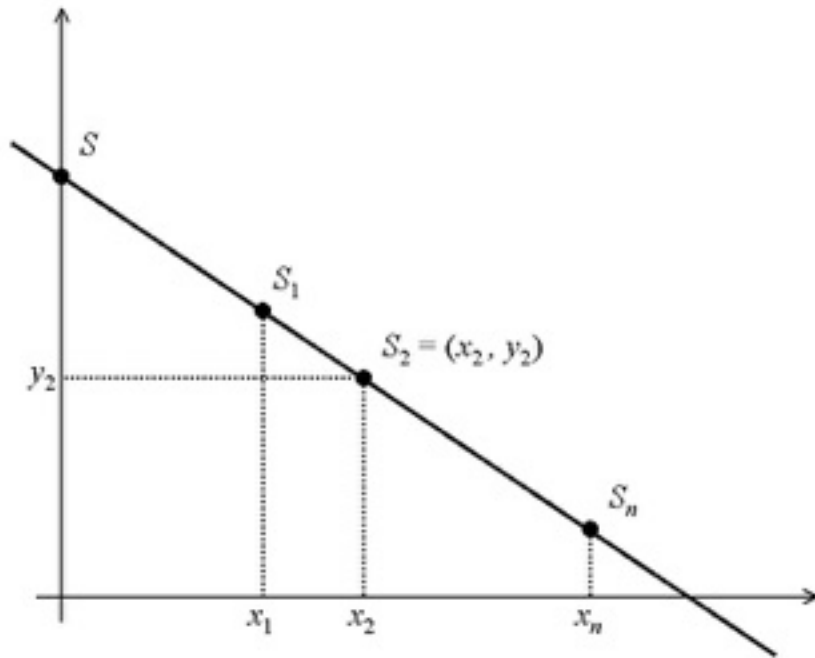
- ◆ Shamir's secret sharing based on polynomial evaluations and Lagrange interpolation is the standard SS scheme
- ◆ To share a secret $s \in F_p$ along $n < p$ players, the dealer D
 - chooses a degree- t polynomial $f_D(x) = s + r_1x + r_2x^2 + \dots + r_tx^t$, for coefficients $r_i \in F_p$ chosen uniformly at random
 - computes and send $y_i = f(i)$ to i^{th} node/party; share $s_i = (i, y_i)$
 - We represent the output of sharing phase as $[s]$
- ◆ Any subset Q of $t+1$ or more players, can reconstruct the secret s as $s = \sum_{i \in Q} \lambda_i y_i$, where $\lambda_i = \prod_{j \in Q \setminus \{i\}} \frac{j}{j-i}$
- ◆ This secret sharing is additive homomorphic

Example: (4,1)-Secret Sharing

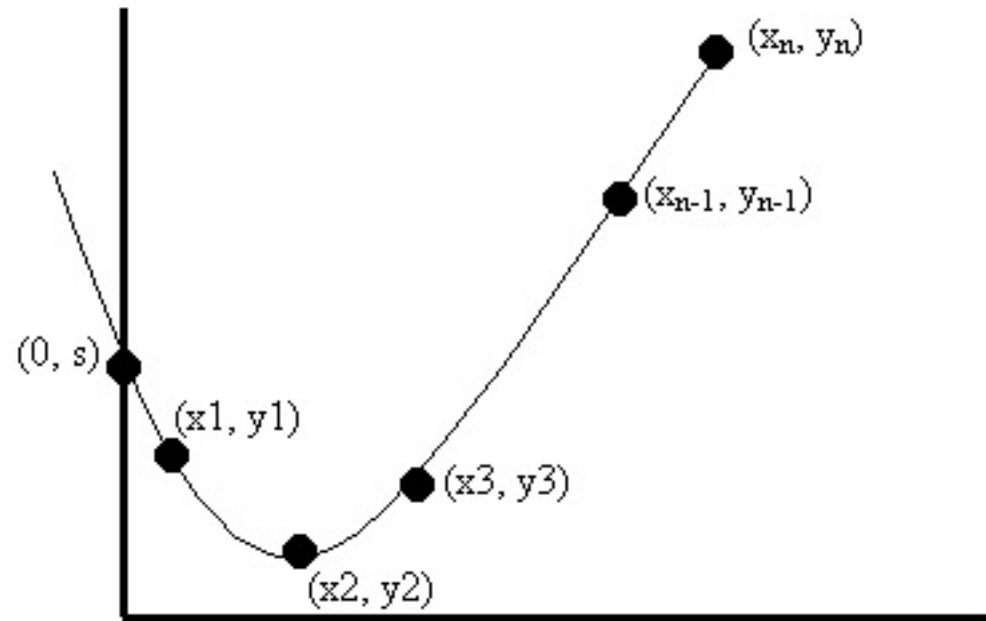


Shamir Secret Sharing Examples

◆ (n,1)-secret sharing

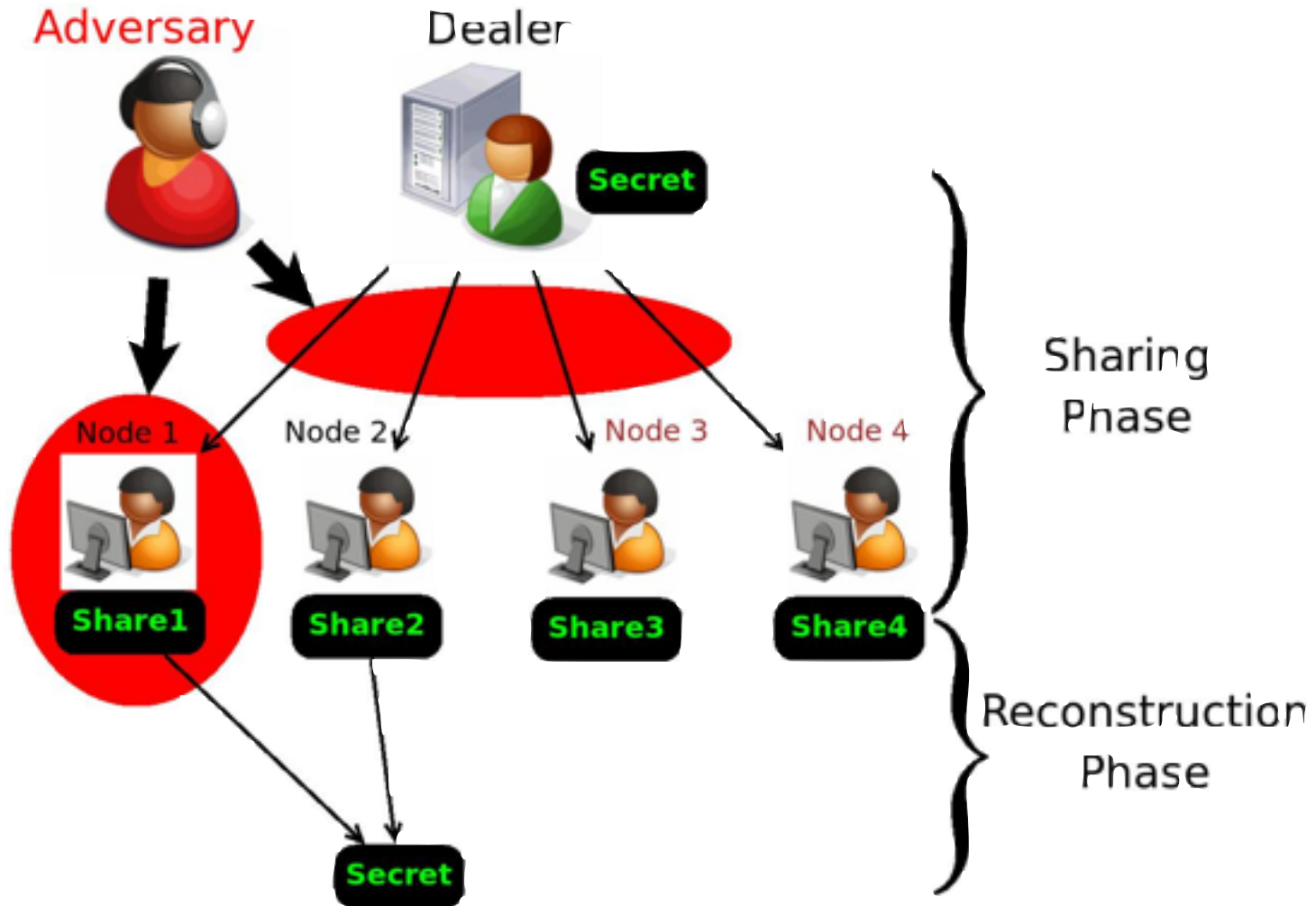


◆ (n,2)-secret sharing



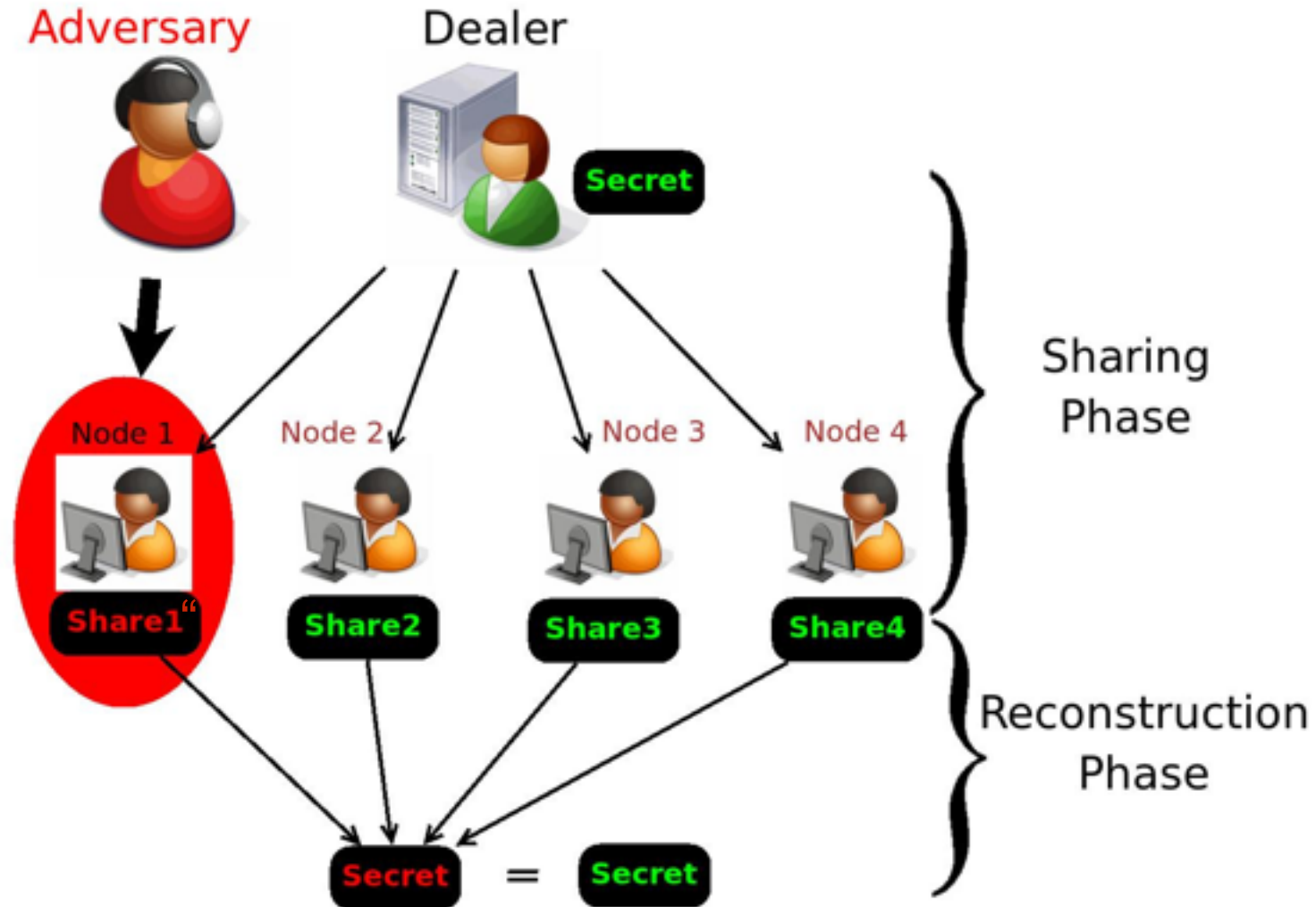
Attacks on Secret Sharing

◆ Attacking Secrecy/Privacy



Attacks on Secret Sharing

◆ Attacking Correctness



Bounding Adversary Behavior

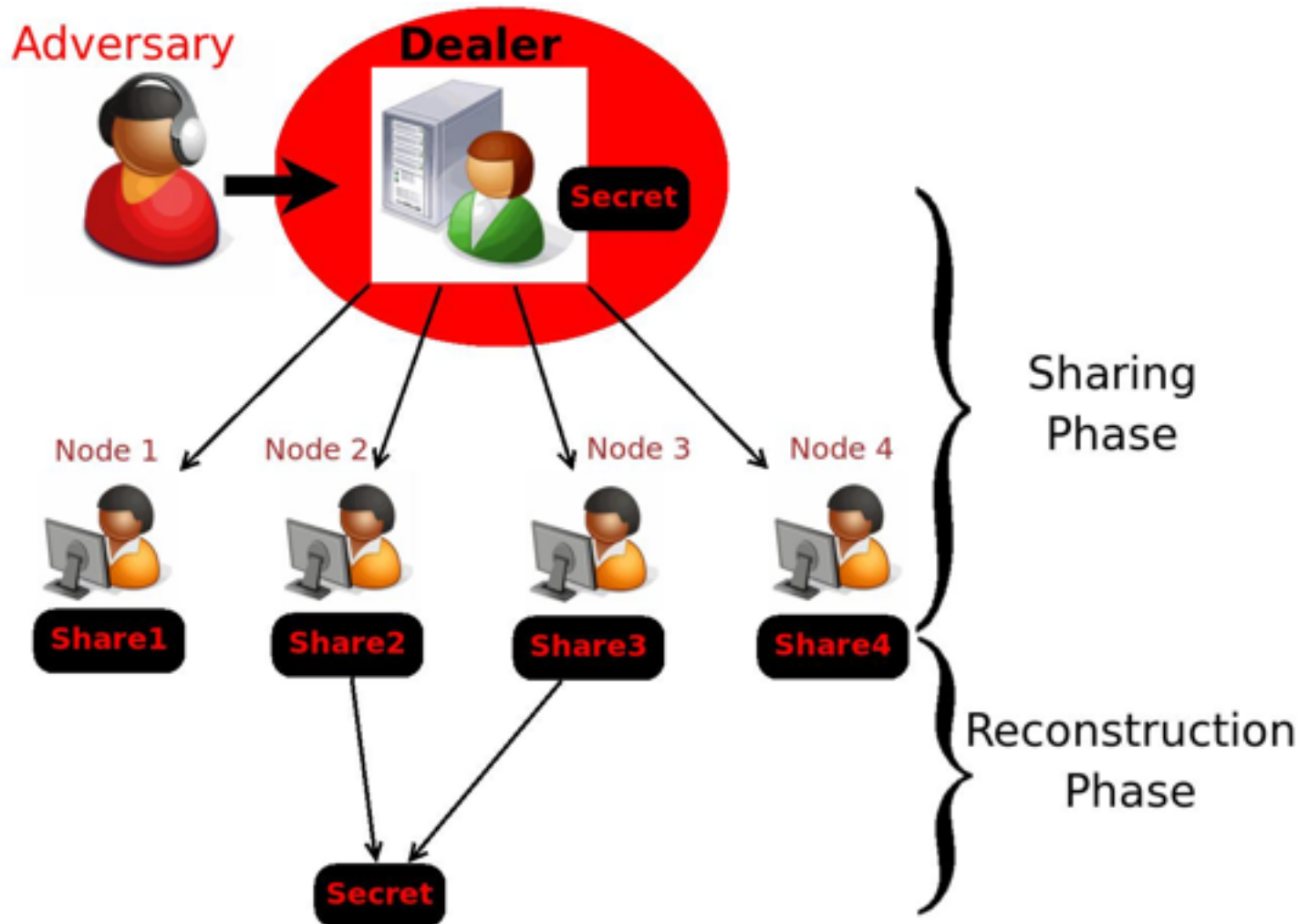
- ◆ What is the relation between n and t
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(Discussion on the board)

- ◆ What if the dealer is malicious?
 - Example Scenarios: Multi-party Computations; Threshold Cryptography

Attacks on Secret Sharing

◆ Attacking Commitment



Verifiable Secret Sharing

- ◆ Secret sharing with three properties
 - Secrecy
 - Correctness
 - Commitment

(Discussion on the board)